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Public-Key Encryption.

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“Period” divides $p-1$.

Today.

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Polynomials.

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Secret Sharing.

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Correcting for loss or even corruption.

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Share secret among n people.

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Two points make a line.

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The idea of the day.

Two points make a line.

Lots of lines go through one point.

Polynomials

A **polynomial**

$$P(x) = a_d x^d + a_{d-1} x^{d-1} \dots + a_0.$$

is specified by **coefficients** a_d, \dots, a_0 .

¹A field is a set of elements with addition and multiplication operations, with inverses. $GF(p) = (\{0, \dots, p-1\}, + \pmod{p}, * \pmod{p})$.

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Note: Often polynomial of degree d means polynomial of at most d .

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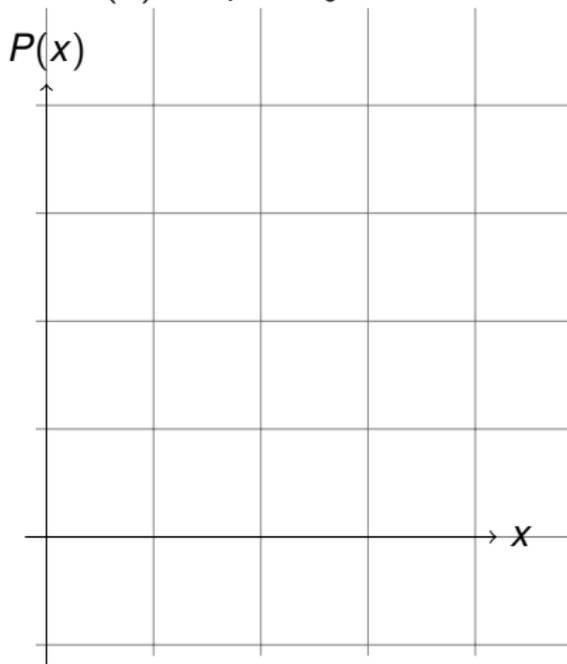
Line: $P(x) = a_1 x + a_0$

Polynomial: $P(x) = a_d x^d + \dots + a_0$

Line: $P(x) = a_1 x + a_0 = mx + b$

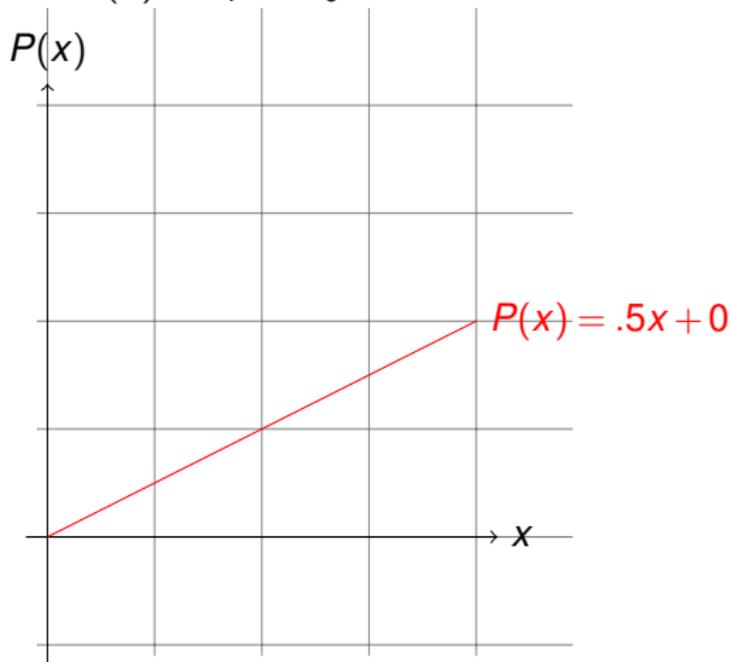
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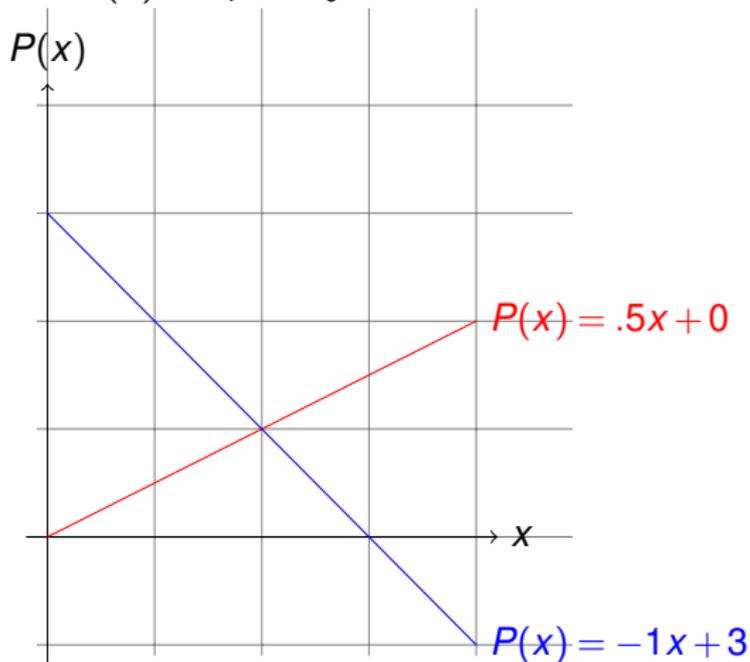
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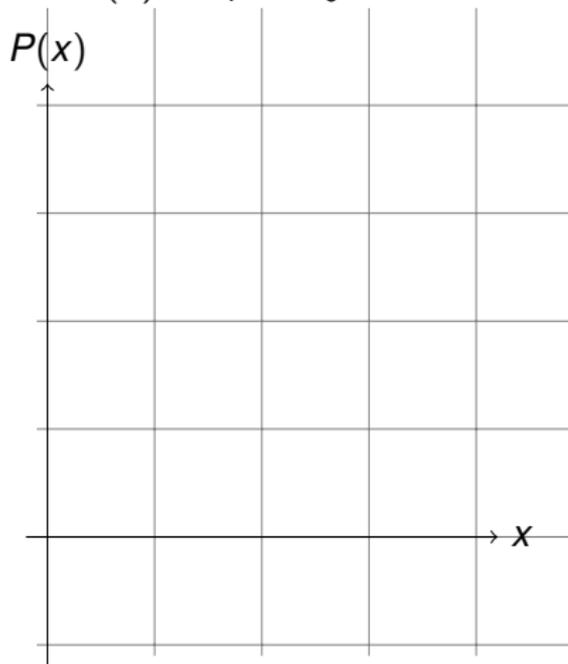
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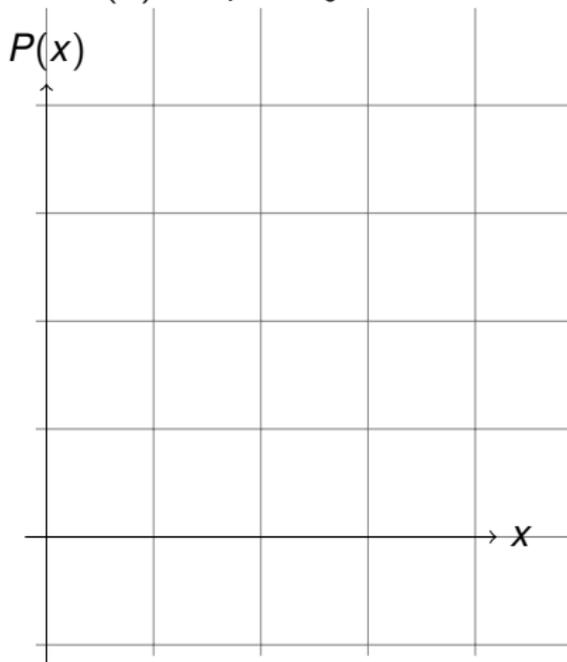
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Parabola: $P(x) = a_2 x^2 + a_1 x + a_0$

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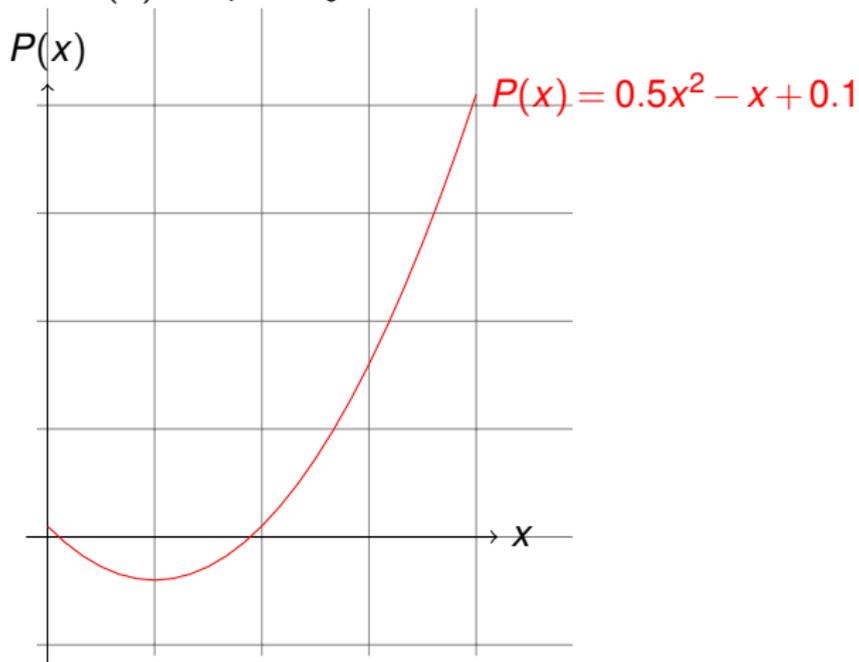
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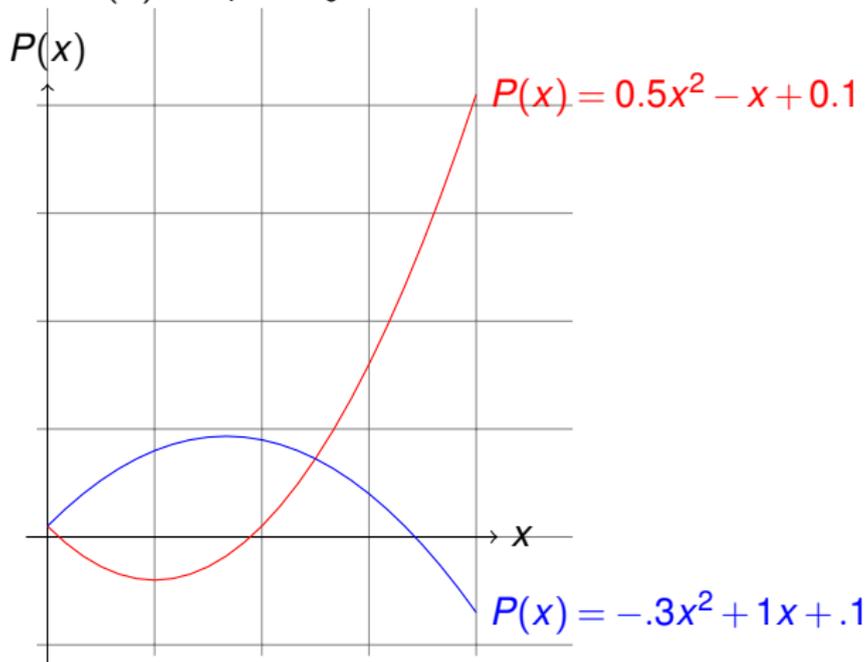
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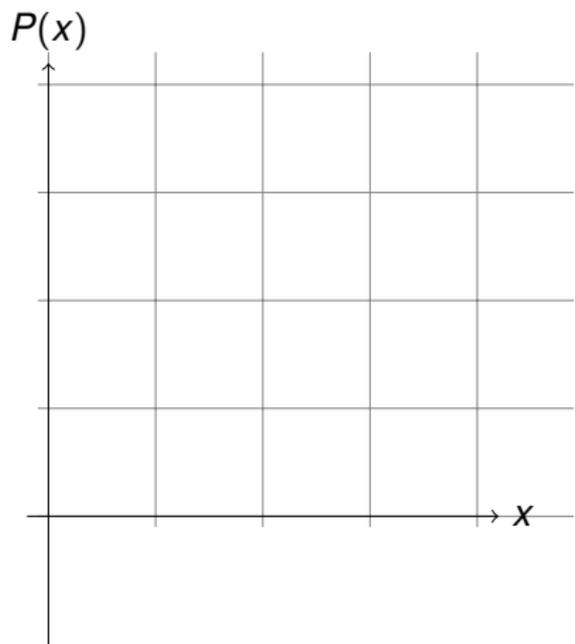
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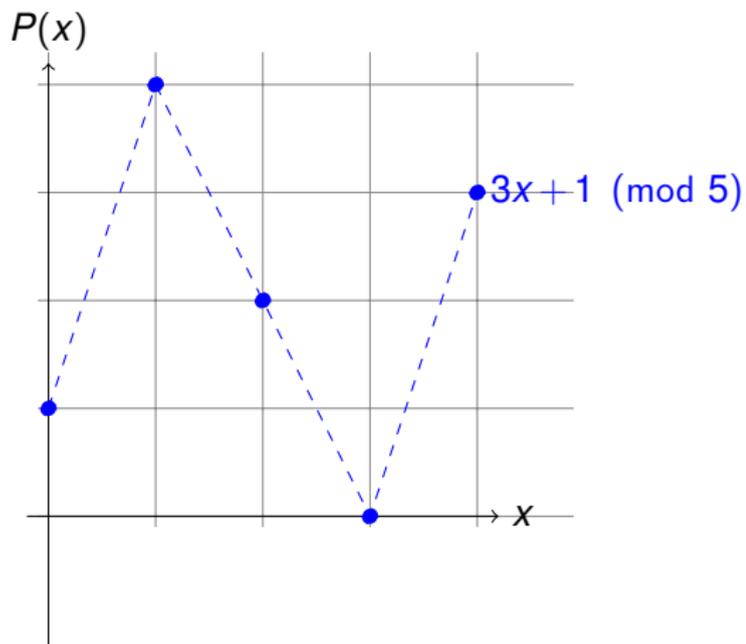


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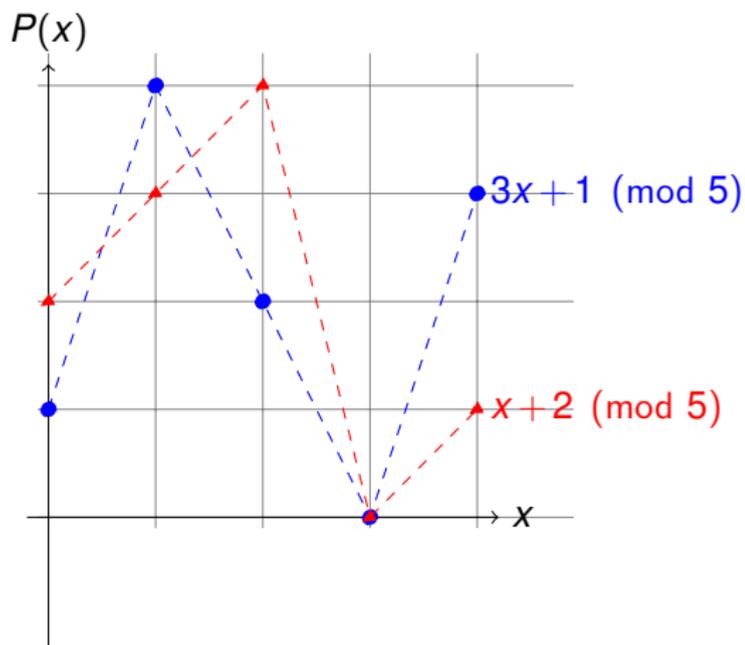
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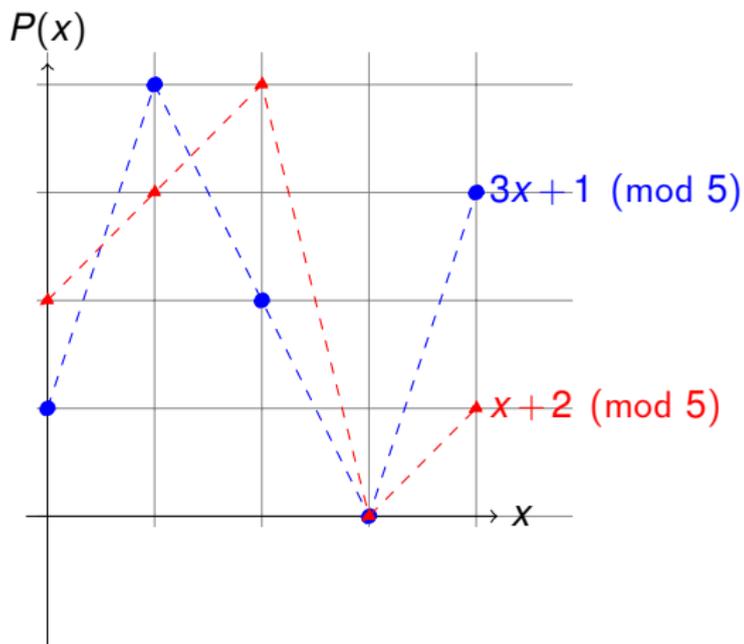


Finding an intersection.

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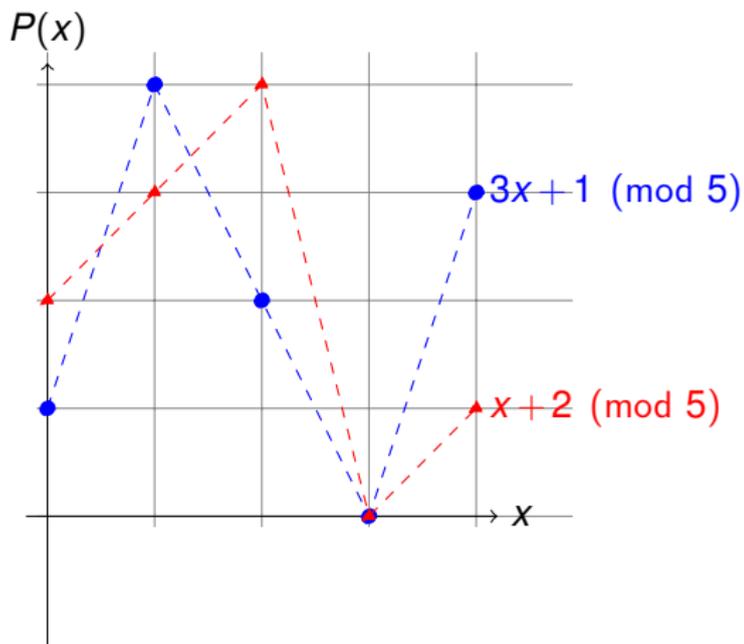
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Good when modulus is prime!!

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Fact: Exactly 1 degree $\leq d$ polynomial contains $d + 1$ points. ²

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Say points are $(x_1, y_1), (x_2, y_2)$.

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All true.

In the Flow (Steph Curry) Poll.

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Flow poll. (All true. (B) is not a proof, it is restatement.)

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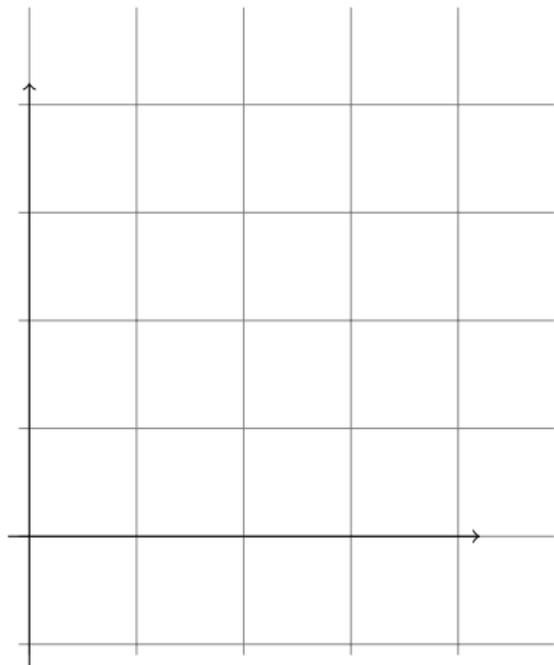
(B) $a_1 = b$

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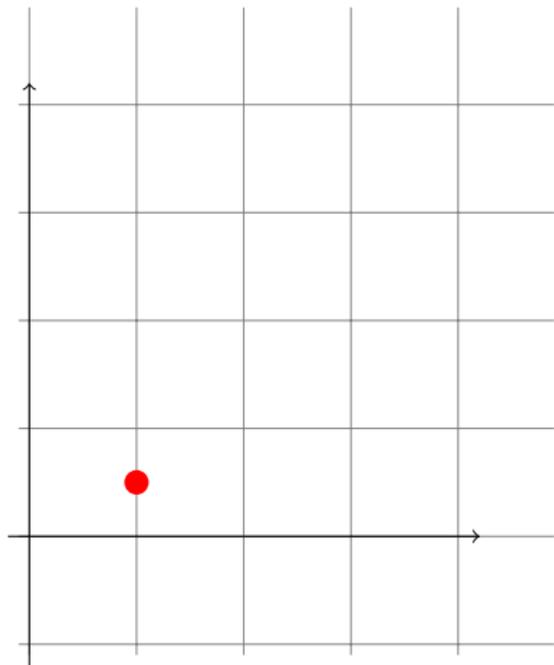
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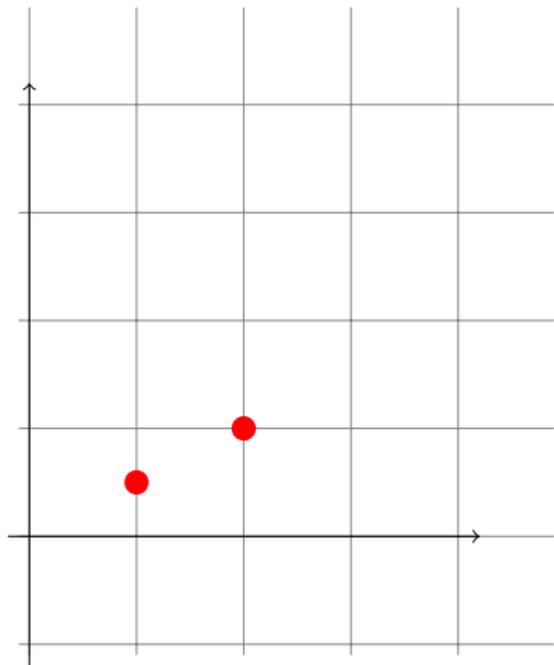
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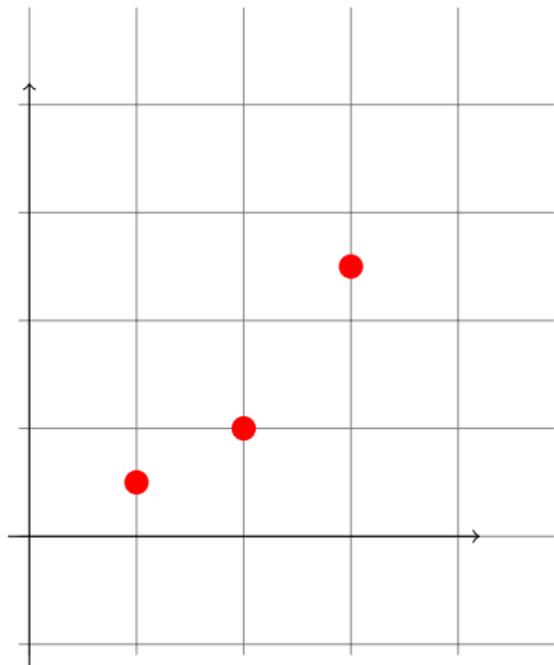
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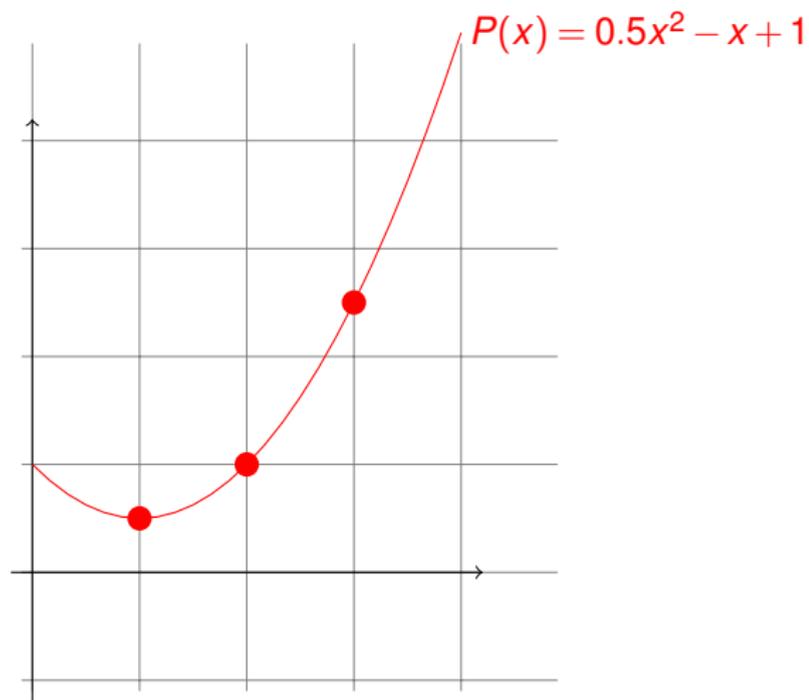
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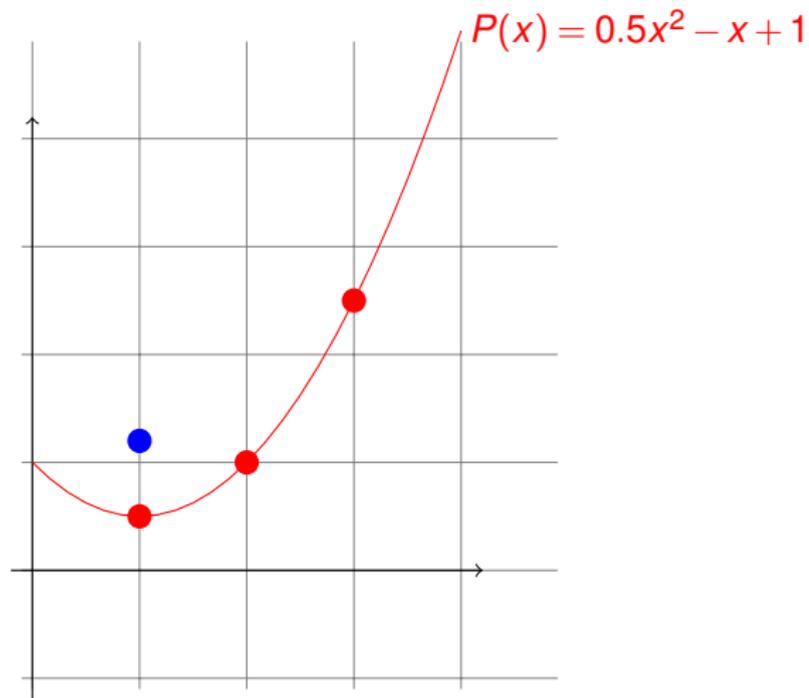
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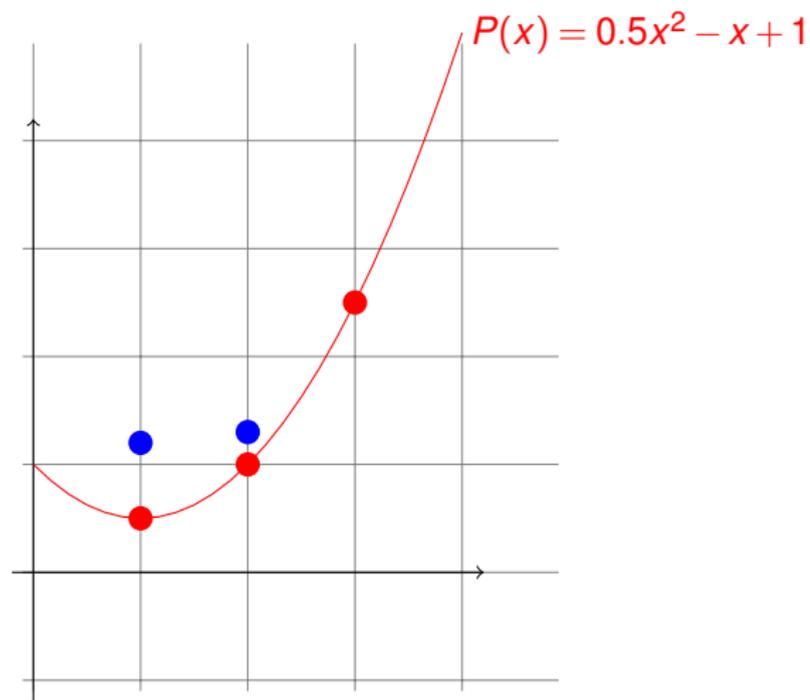
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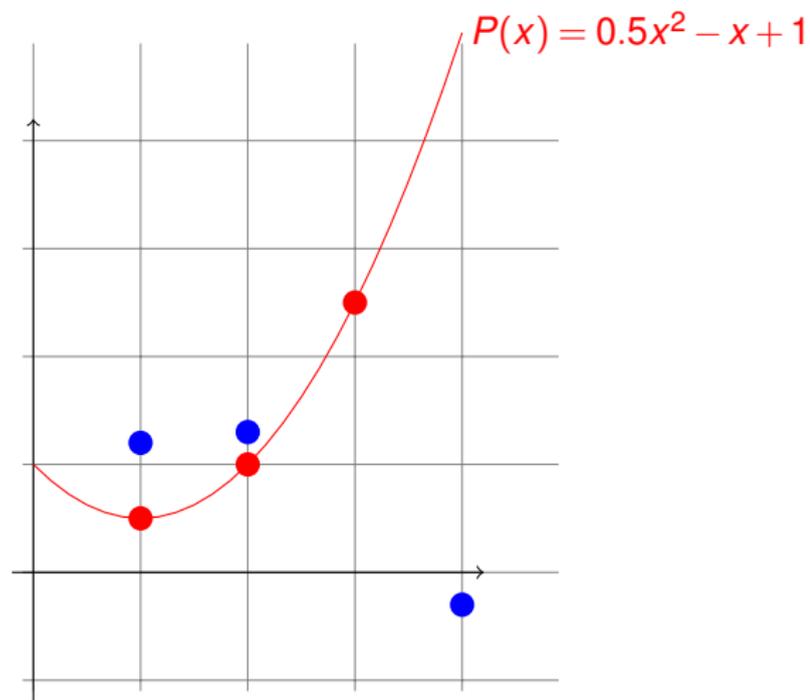
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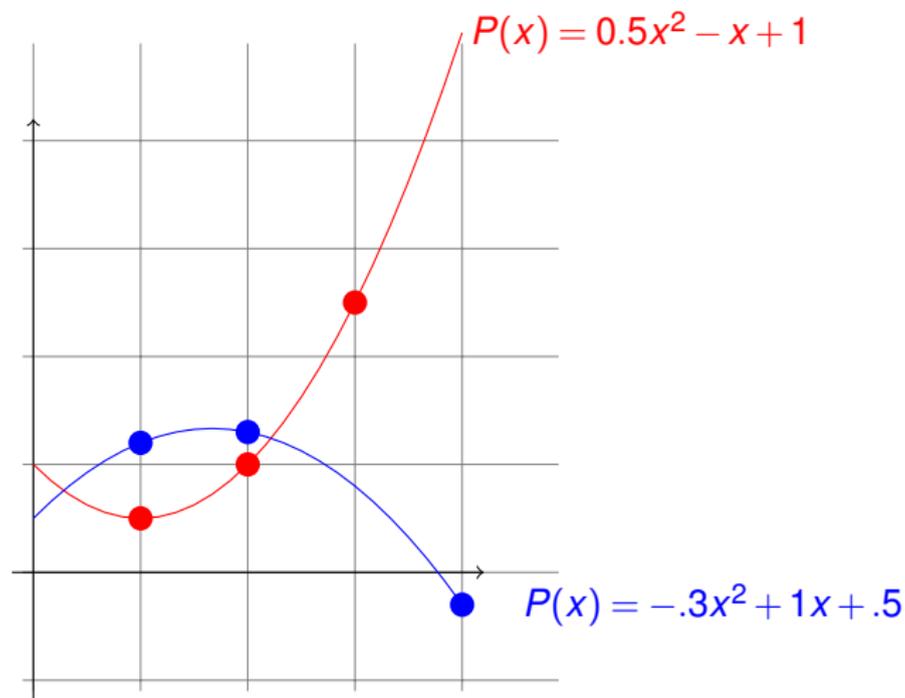
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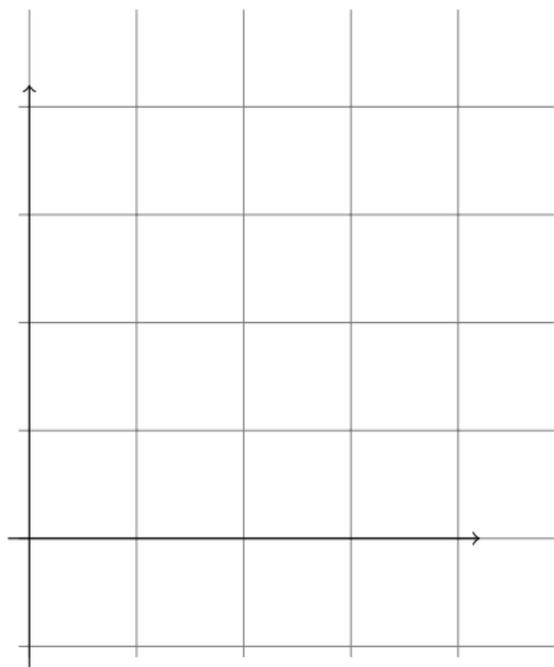
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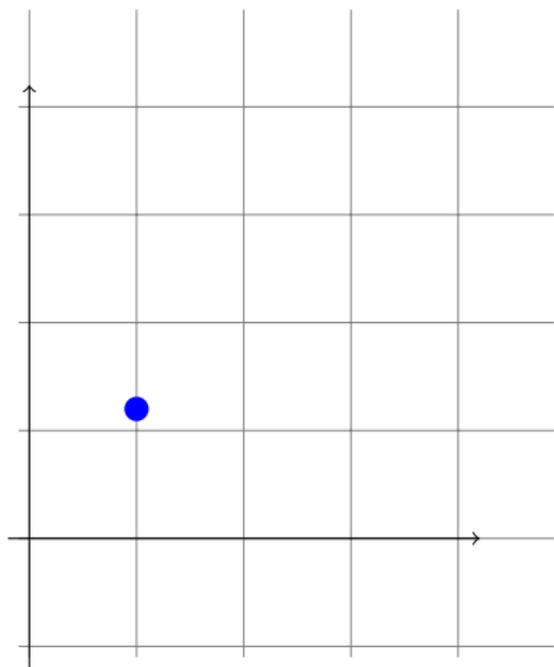
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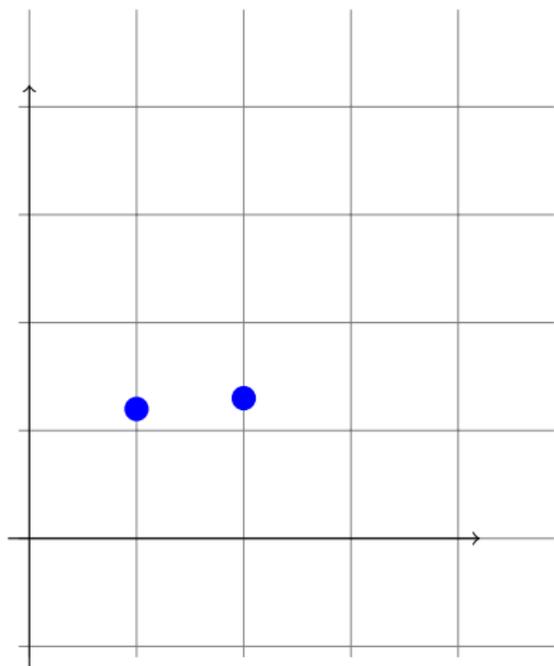
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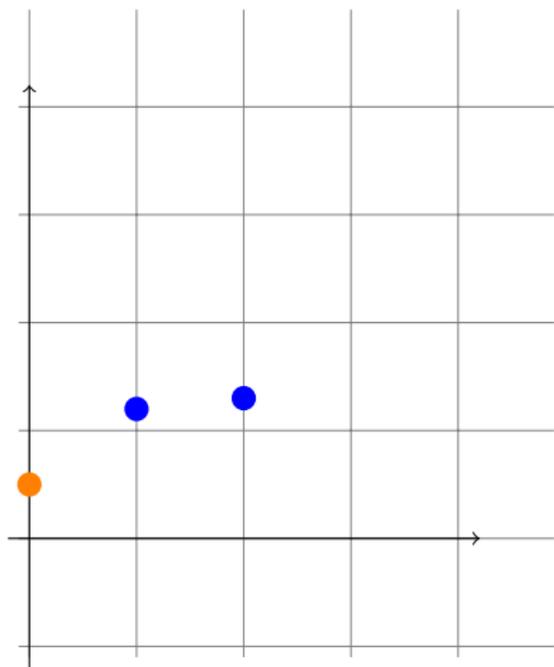
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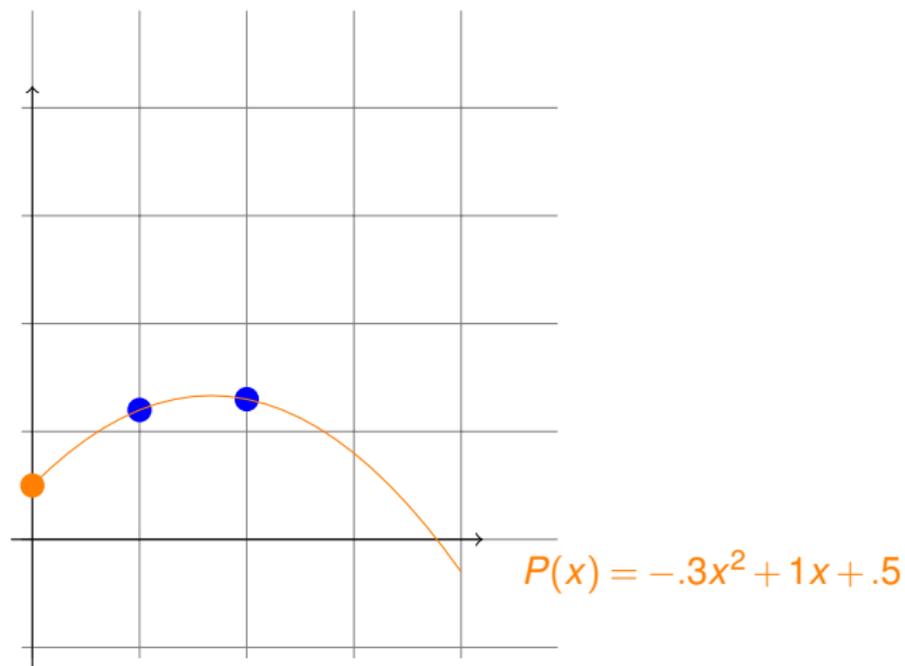
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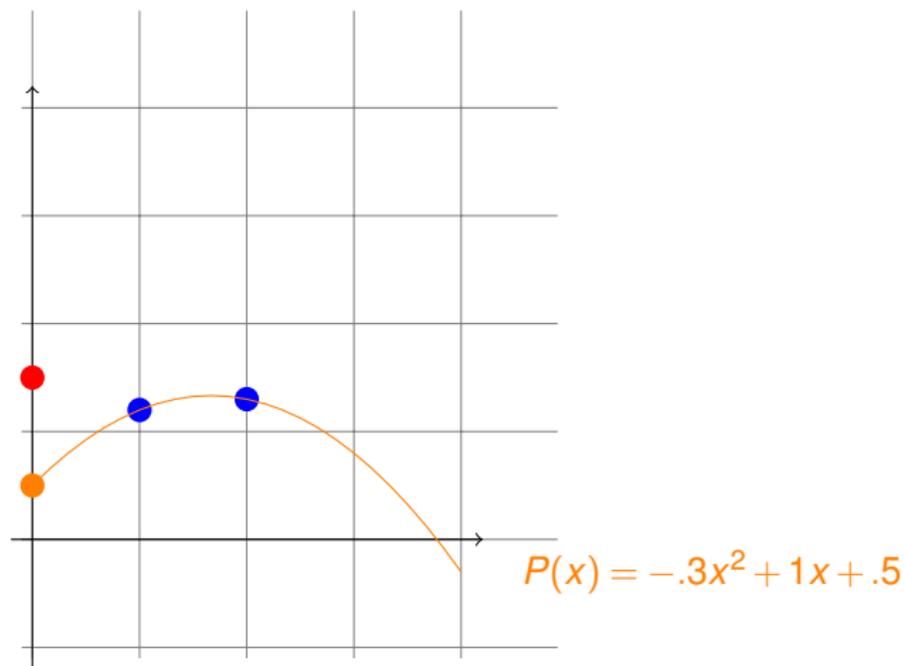
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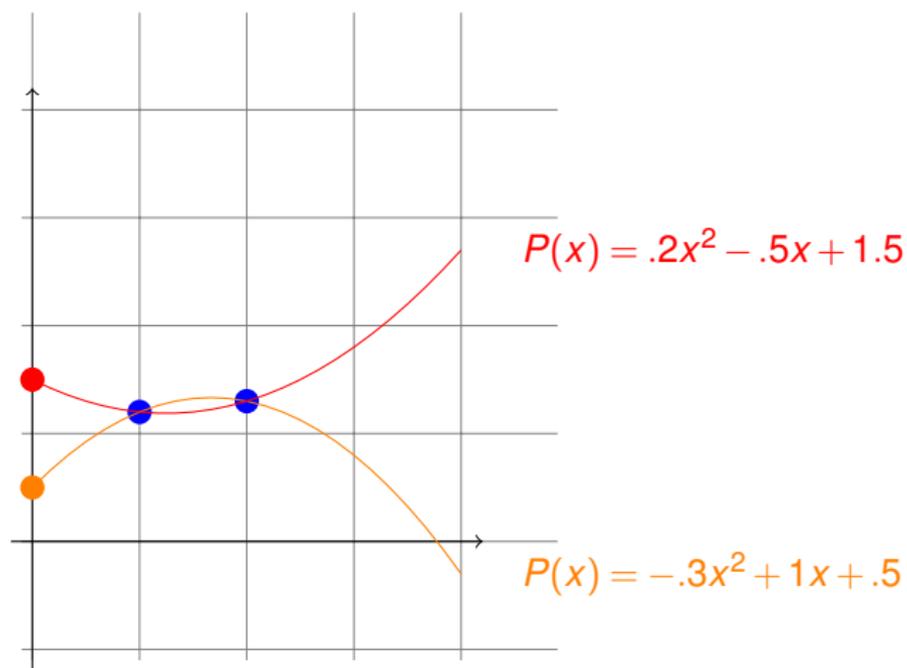
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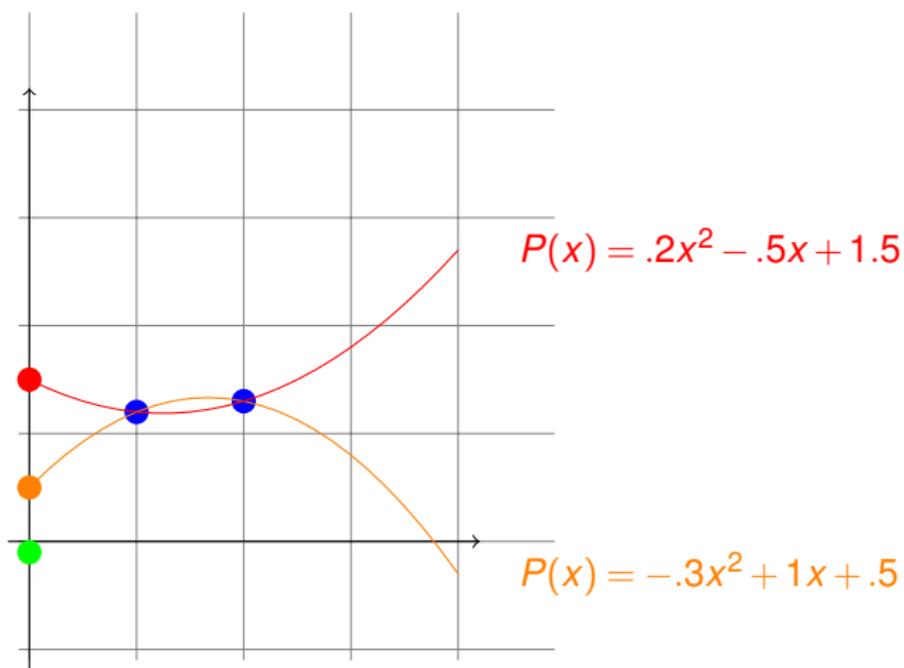
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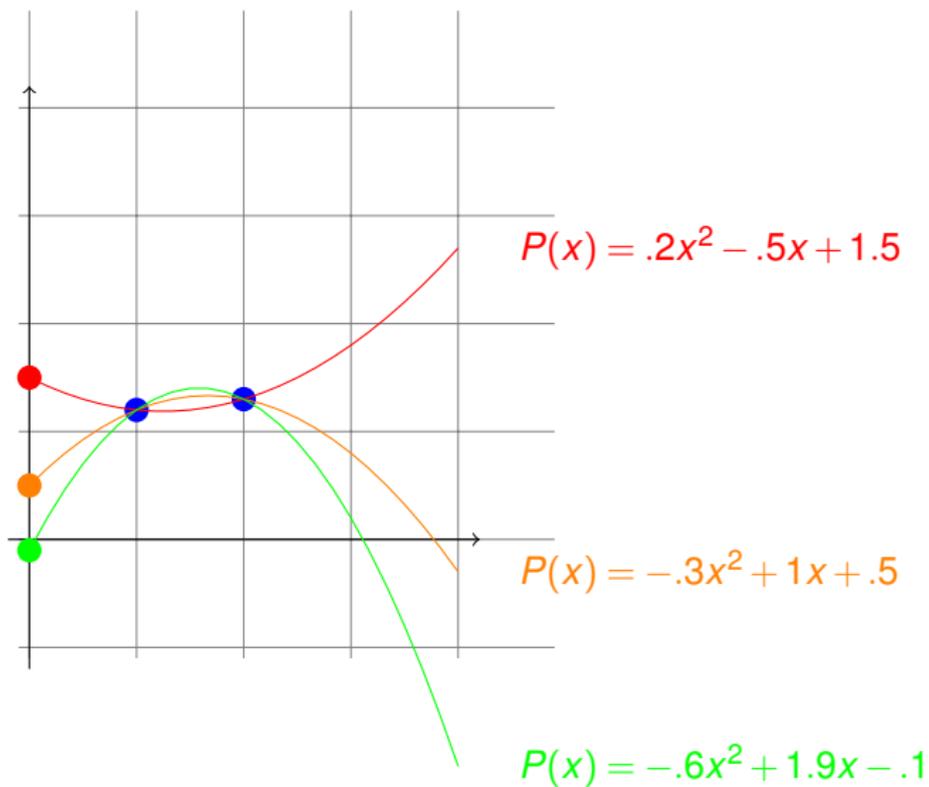
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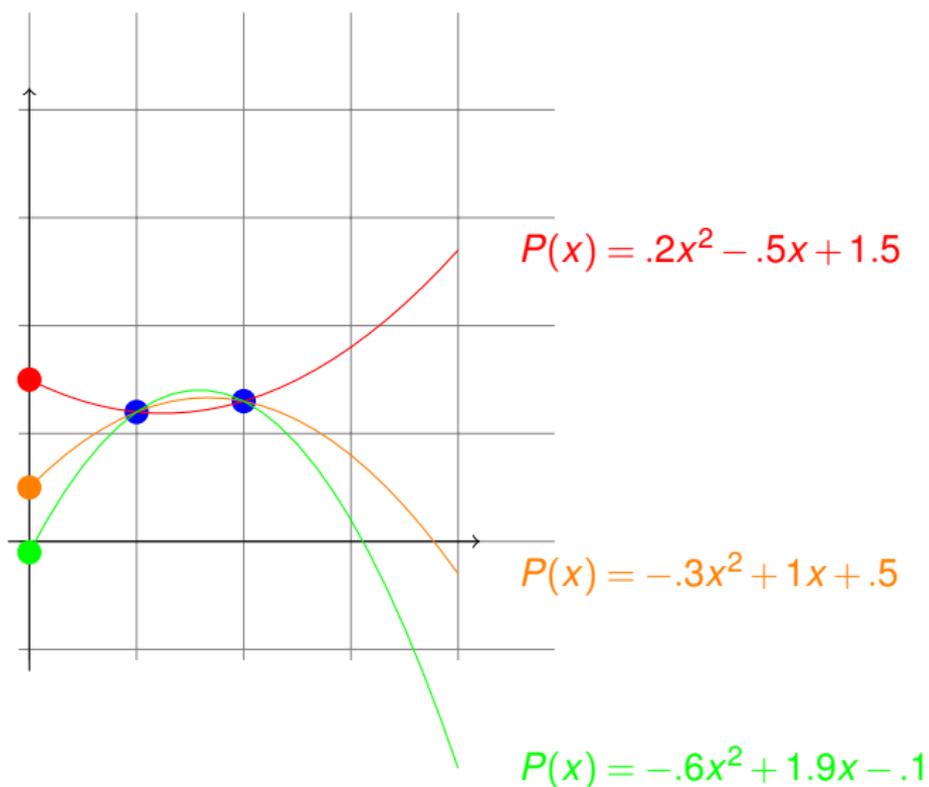


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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

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Given points: $(x_1, y_1); (x_2, y_2) \cdots (x_k, y_k)$.

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Numerator is 0 at $x_j \neq x_i$. "Denominator" makes it 1 at x_i .

$$\Delta_i(x_j) = 0 \text{ if } i \neq j \text{ and } \Delta_i(x_i) = 1$$

And..

$$P(x) = y_1 \Delta_1(x) + y_2 \Delta_2(x) + \cdots + y_{d+1} \Delta_{d+1}(x).$$

hits points $(x_1, y_1); (x_2, y_2) \cdots (x_{d+1}, y_{d+1})$.

$$\text{Since } P(x_i) = y_1(0) + y_2(0) \cdots + y_i(1) \cdots + y_{d+1}(0).$$

And Degree d polynomial.

Construction proves the existence of a polynomial!

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(D) $\Delta_1(x_3) = 1$

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(B), (C), and (E)

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Put the delta functions together.

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Must prove **Roots fact**.

Polynomial Division.

Divide $4x^2 - 3x + 2$ by $(x - 3)$ modulo 5.

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Proof Sketch: By induction.

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Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.

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Modular Arithmetic Fact: Exactly one polynomial degree $\leq d$ over $GF(p)$, $P(x)$, that hits $d + 1$ points.

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(Almost) the same as what is missing: one $P(i)$.

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1. Evaluate degree $k - 1$ polynomial n times using $\log p$ -bit numbers.
2. Reconstruct secret by solving system of k equations using $\log p$ -bit arithmetic.

A bit more counting.

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Infinite number for reals, rationals, complex numbers!

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Can hand out n points on polynomial as shares.